



Rent's Rule and Parallel Programs: Characterizing Network Traffic Behavior

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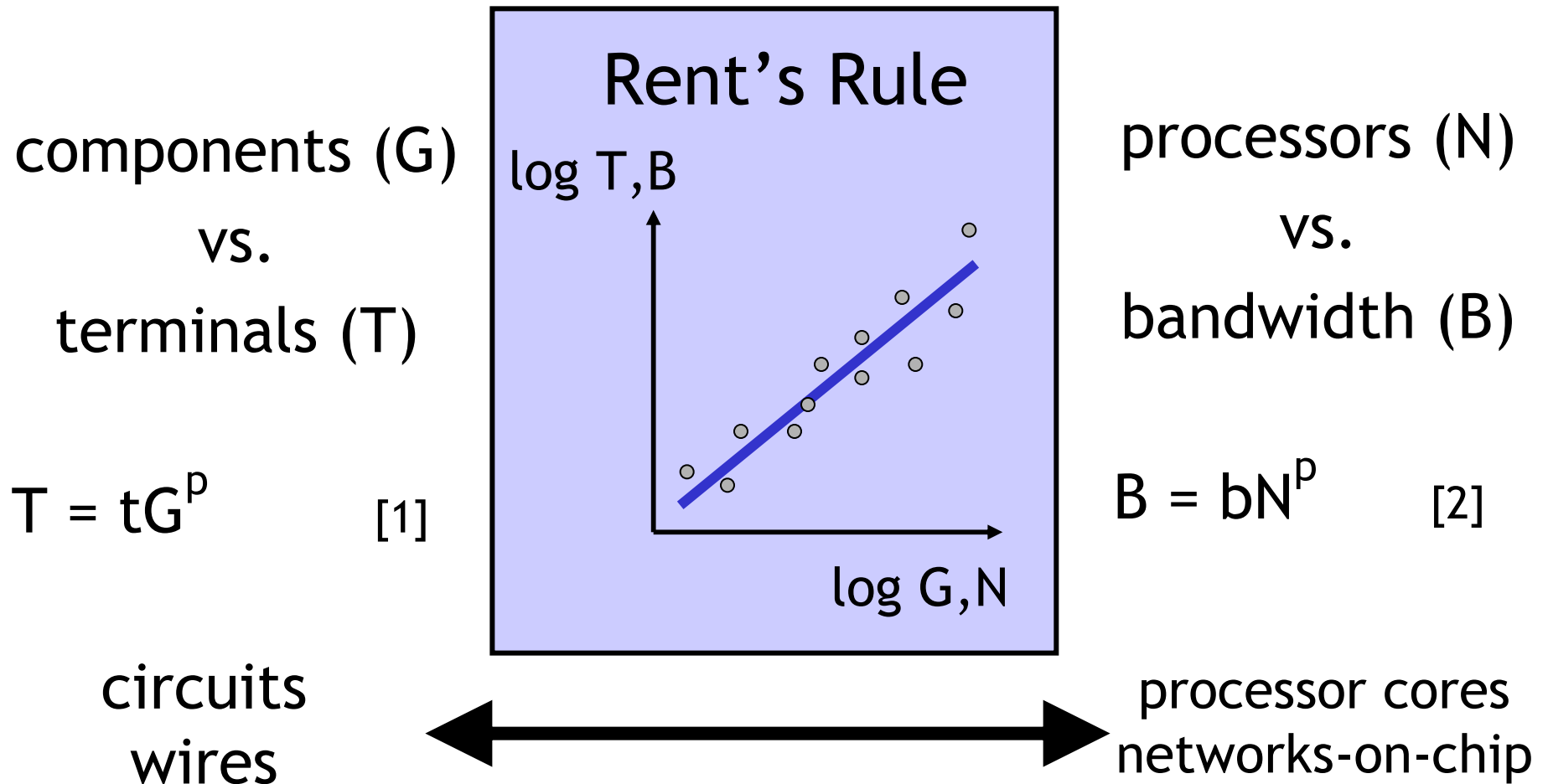
Outline

- Introduction
- Rent's rule & traffic locality
- Time-varying network traffic
- Conclusions

Evolution of Systems design

- VLSI systems get ever more complicated
- More software, processor IP blocks, hardware/software co-design
- Ad-hoc global wiring → Network-on-Chip (“communication IP block”), long wires → packets
- What with Rent’s rule?

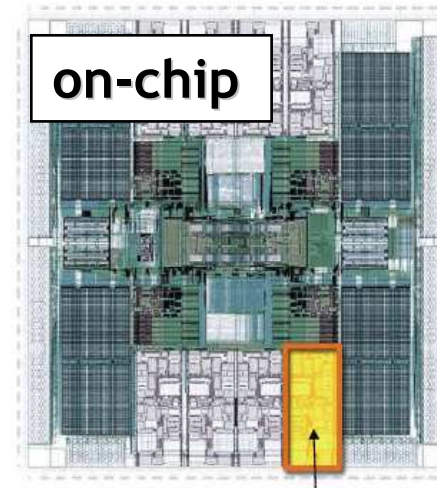
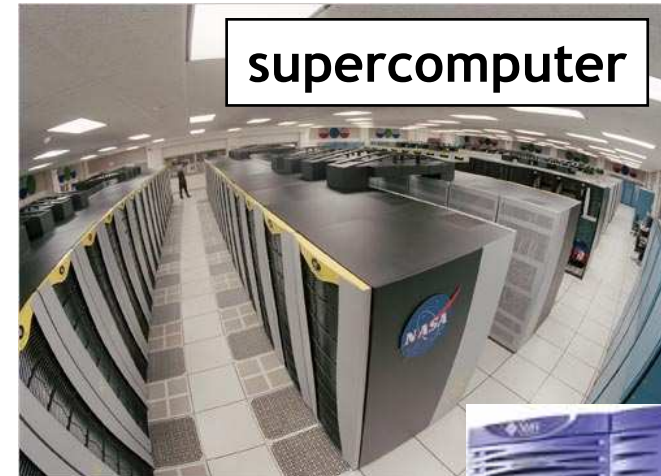
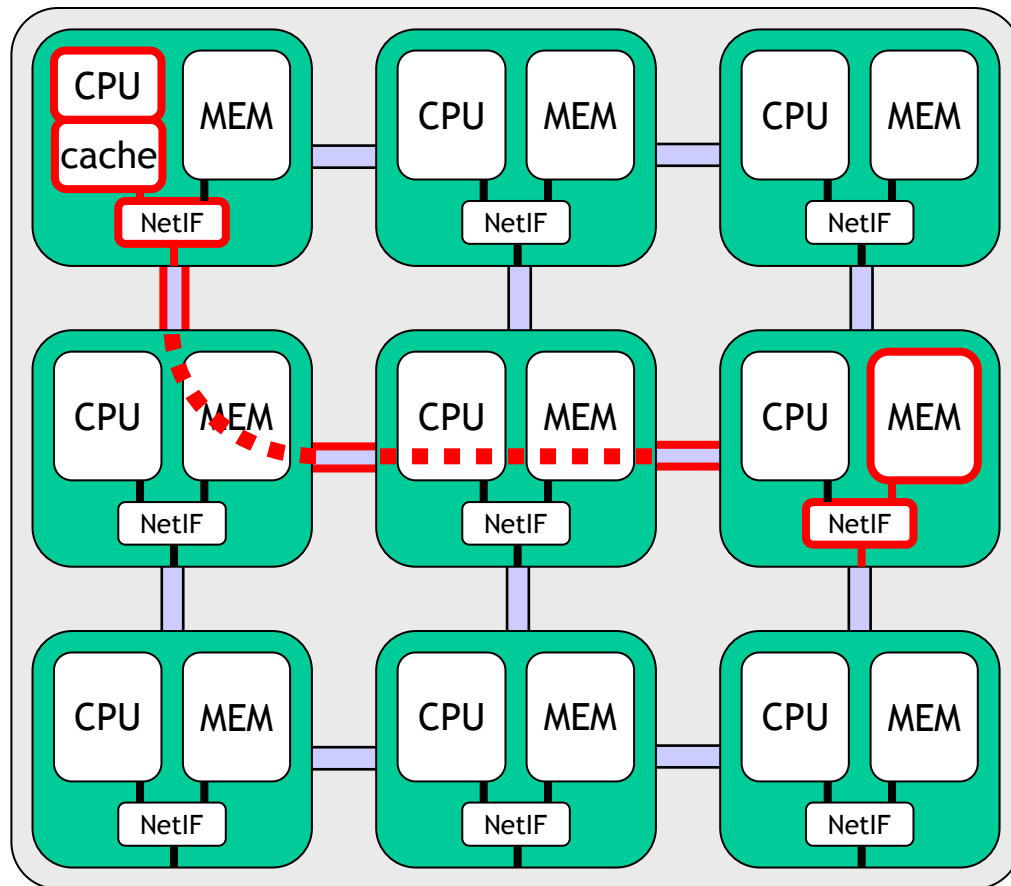
Rent's rule: power law relation



- [1] Landman and Russo, IEEE Trans. on Computers, 1971
[2] D. Greenfield et. al, NOCS 2007.

Multiprocessor + Network architecture

Shared memory: network is part of memory hierarchy



UltraSPARC-Core

NoC design: problems and opportunities

- Simple traffic models: uniform, hot-spot, fixed bandwidth distribution
 - Ignores locality, time-variance in network traffic
 - Yields non-optimal NoC designs (uniform vs. non-uniform, static vs. reconfigurable)
- Opportunity: better traffic models, analytical tools vs. trial-and-error

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Partitioning nodes by communication intensity

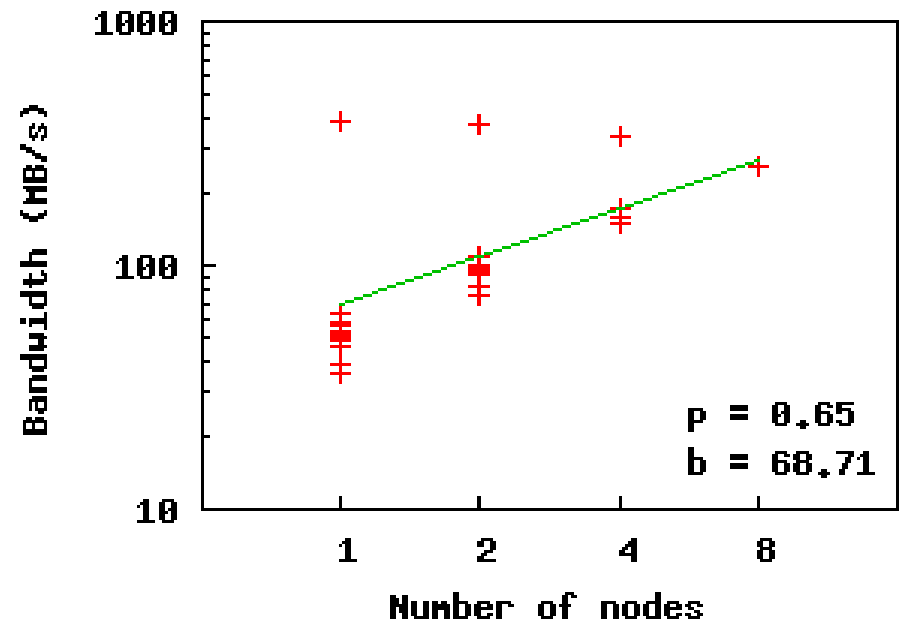
- Hierarchically partition nodes according to communication (hMETIS)
- Just as for wires, but:
 - Communication graph is usually fully connected
 - Weight on each connection
= total communication between node pair
- Fit power law on (cluster size, bandwidth) distribution

Rent exponent

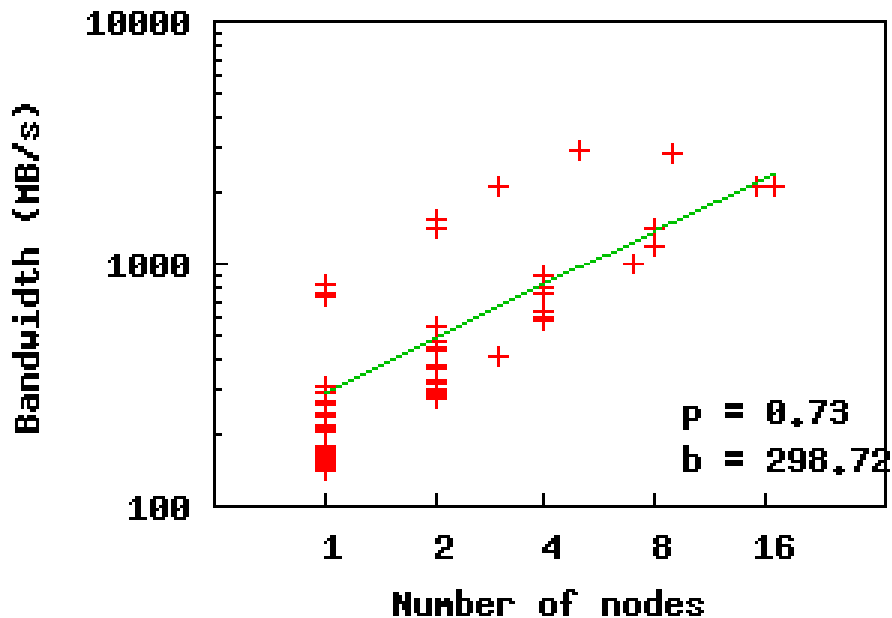
measured Rent
exponent (dependent
on application):

- 16 nodes: .55-.65
- 64 nodes: .66-.74

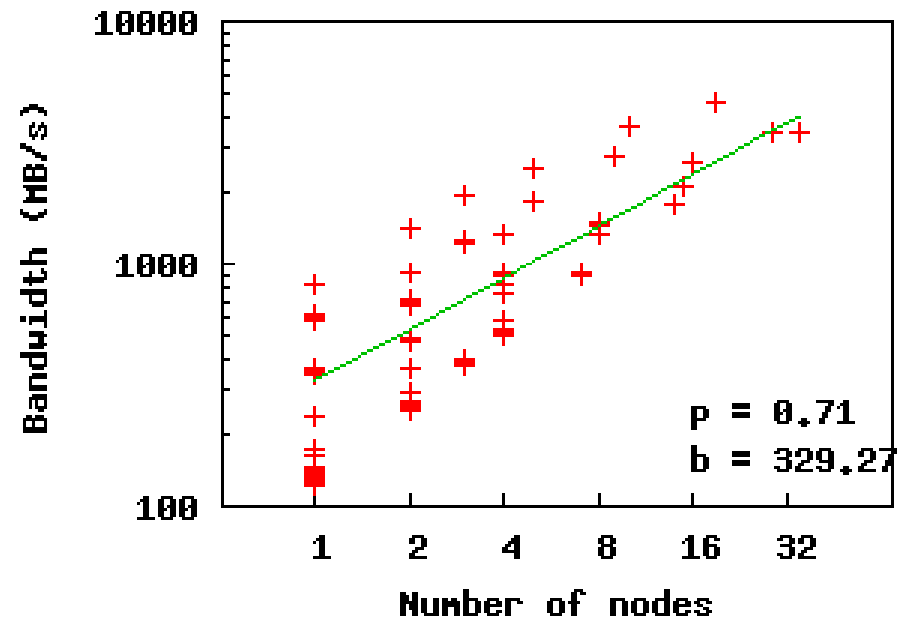
cholesky29, 16 nodes



ocean.cont, 32 nodes



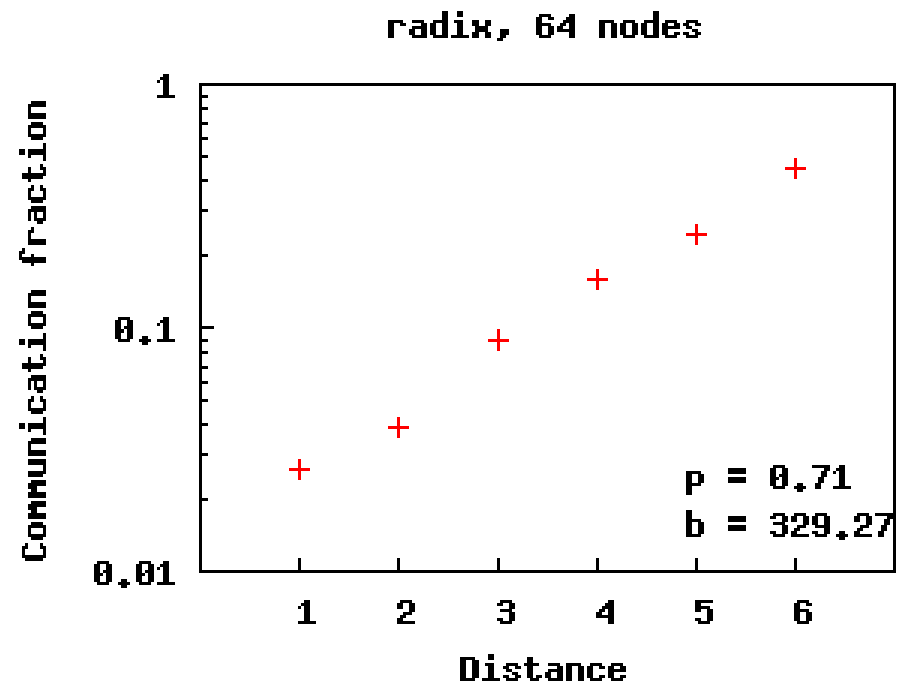
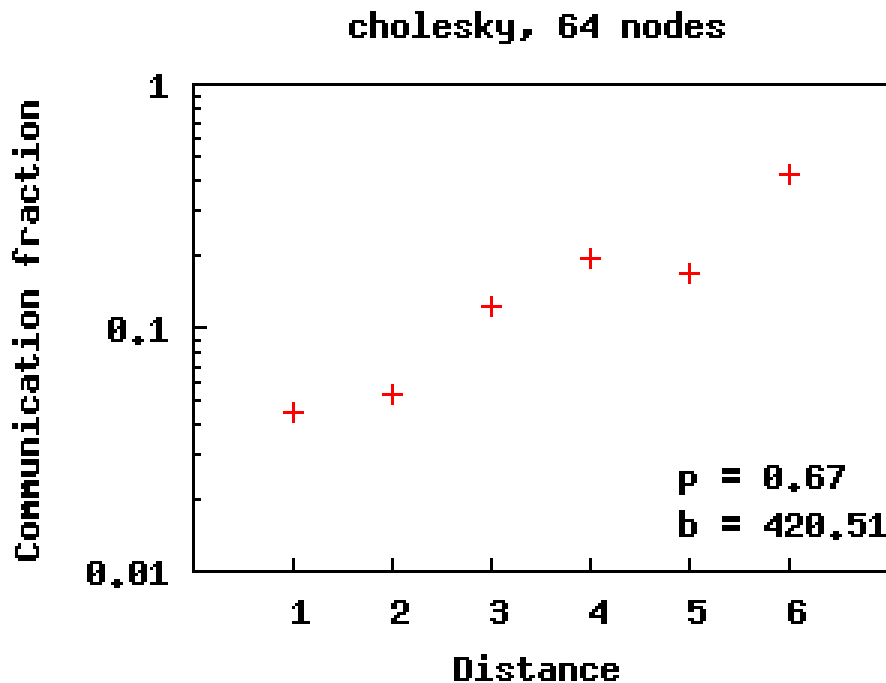
radix, 64 nodes



“Wire length” distribution

Distribution of communication vs. distance

distance(A, B) =
 $\log_2(\text{size of smallest cluster containing both A and B})$



Outline

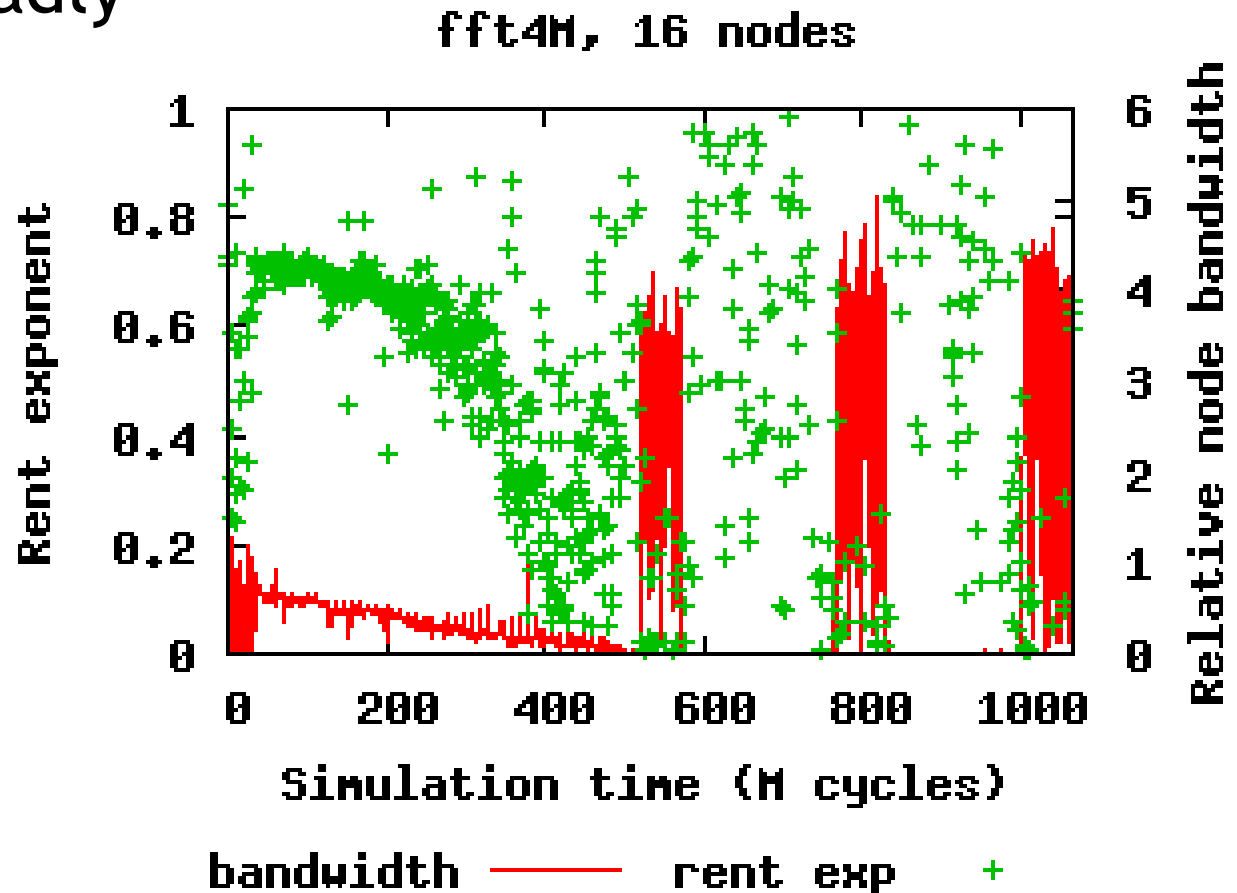
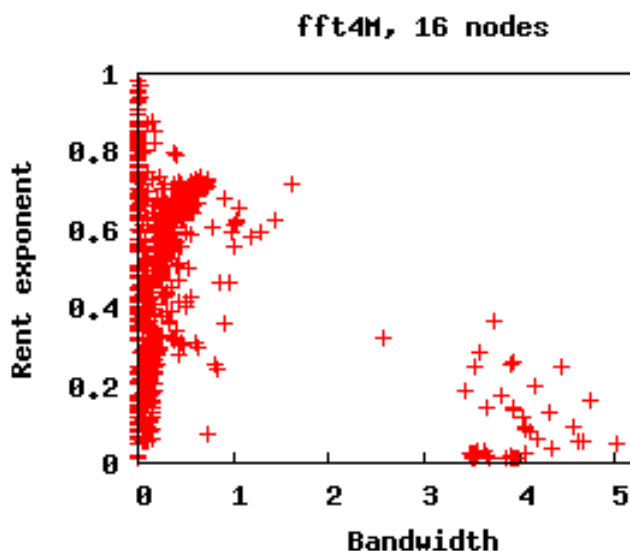
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Communication varies through time

- Hardware:
 - fixed function
 - traffic remains similar through time
- Software:
 - more complex, different phases (e.g. function call)
 - communication patterns can change through time

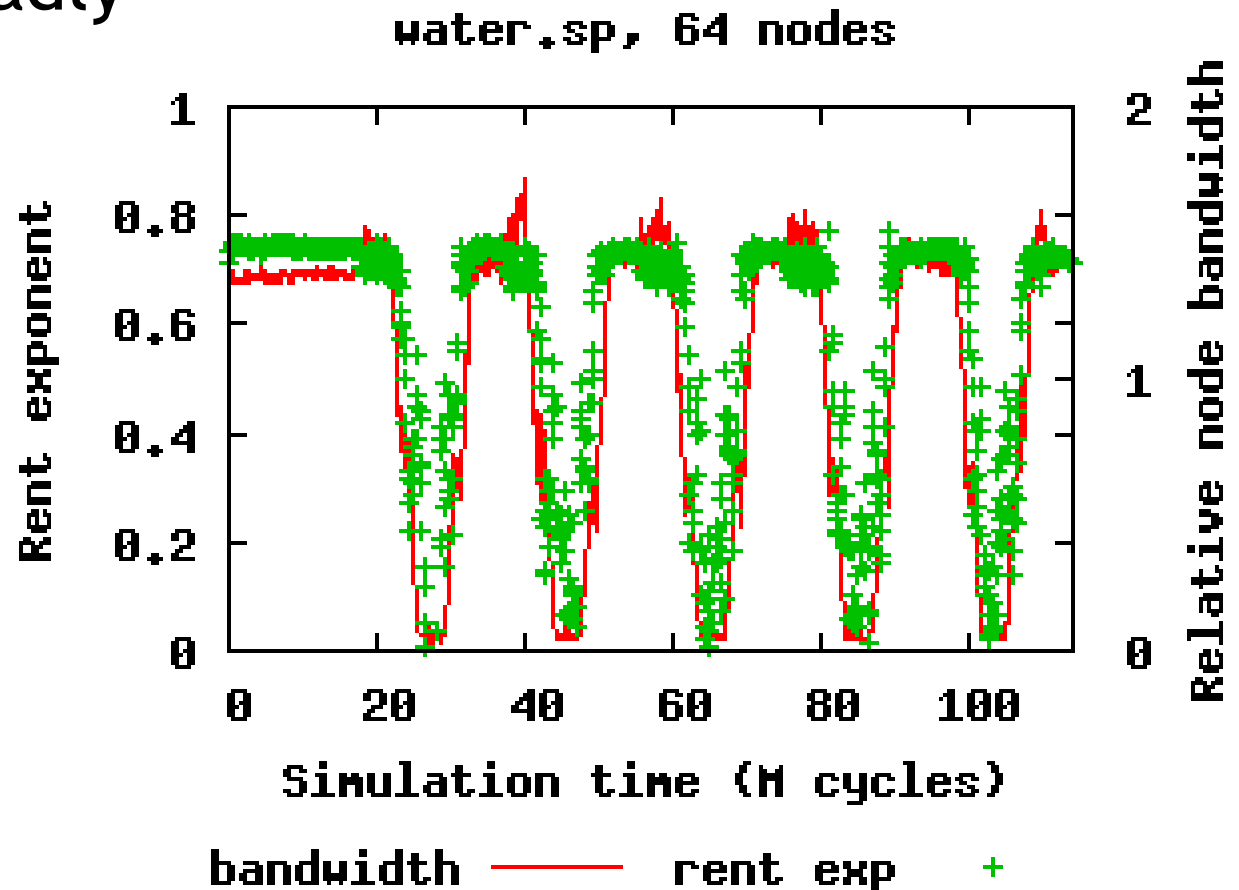
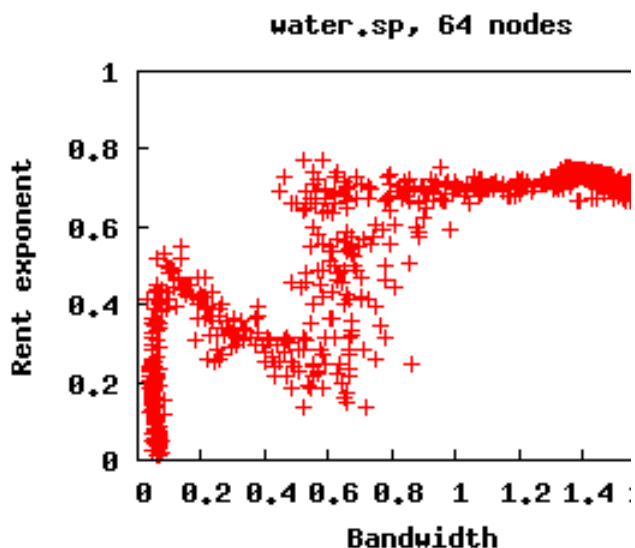
Communication varies through time

- Repeat partitioning per interval of 100k clock cycles
- Periods of high and low communication alternate
- Rent exponent badly defined during periods of low communication



Communication varies through time

- Repeat partitioning per interval of 100k clock cycles
- Periods of high and low communication alternate
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Node placement vs. variable traffic

- Node partitioning can lead to optimal node placement (minimal communication distances)
- But: varying traffic \rightarrow varying optimal placement?
- Compute interval similarity, based on partitionings
- Account for traffic intensity (moving non-communicating nodes has no effect)

Similarity of communication between intervals

- For time intervals X and Y , each with traffic pattern *traffic* and optimal partitioning *part*
- $\text{part}[X]$ cuts minimal fraction of $\text{traffic}[X]$
- assume we use $\text{part}[X]$ in interval Y , what fraction of $\text{traffic}[Y]$ is cut? $\rightarrow \text{cut}[X, Y]$
- always more than $\text{part}[Y]$ would = $\text{cut}[Y, Y]$
- similarity of partitionings, accounting for traffic intensity:

$$\text{sim}[X, Y] = \frac{\text{cut}[X, X] + \text{cut}[Y, Y]}{\text{cut}[Y, X] + \text{cut}[X, Y]}$$

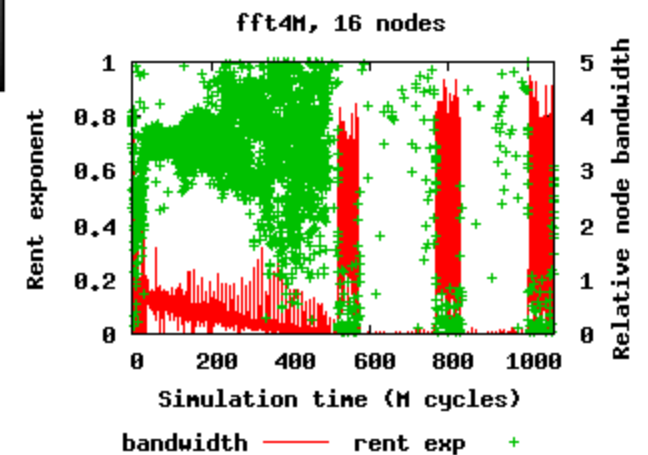
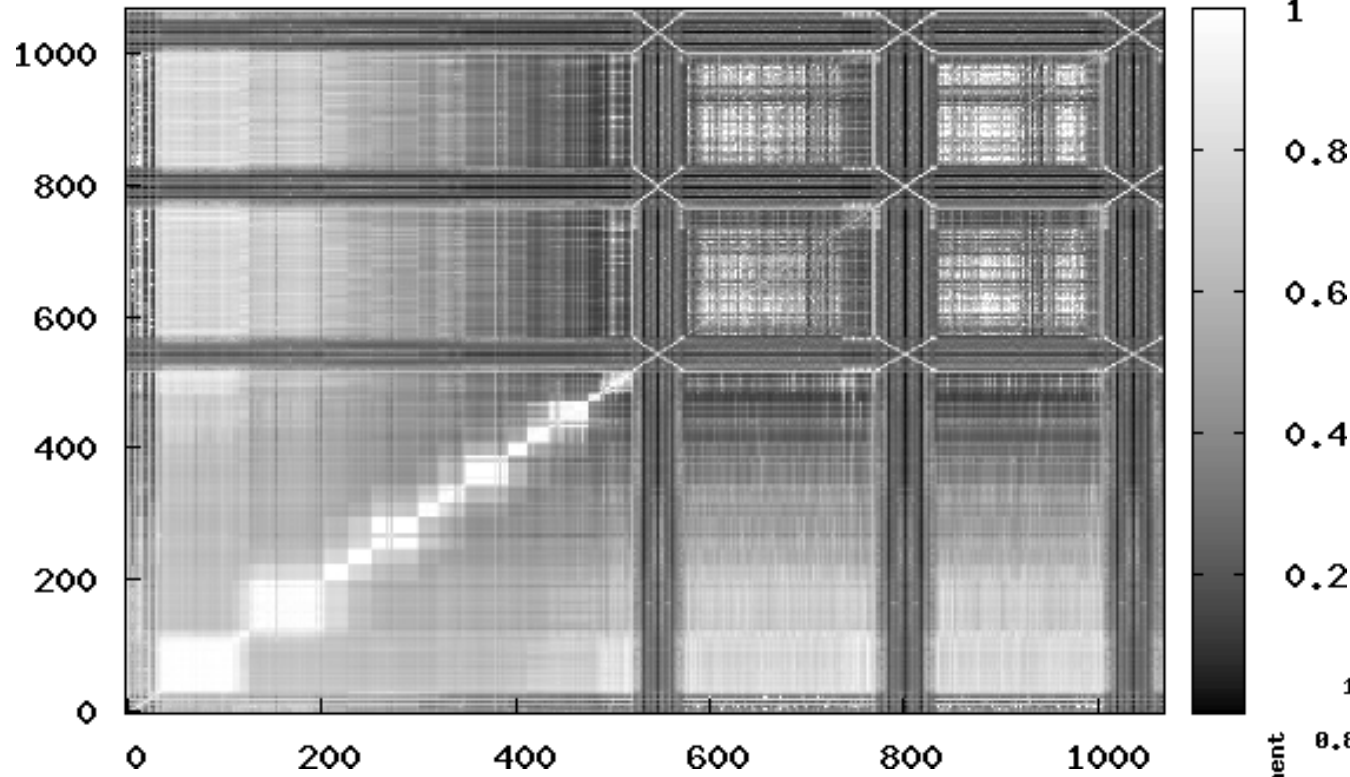
Similarity measure properties

$$\text{sim}[X, Y] = \frac{\text{cut}[X, X] + \text{cut}[Y, Y]}{\text{cut}[Y, X] + \text{cut}[X, Y]}$$

- $\text{cut}[X, X] < \text{cut}[Y, X]$ and $\text{cut}[Y, Y] < \text{cut}[X, Y]$
 $\rightarrow 0 \leq \text{sim}[X, Y] \leq 1$
- $\text{sim}[X, X] = 1$
- when $\text{traffic}[X] \gg \text{traffic}[Y]$:
 $\text{cut}[* , Y] \sim 0 \rightarrow \text{sim}[X, Y] \sim \text{cut}[X, X] / \text{cut}[Y, X]$
(only dependent on $\text{traffic}[X]$)

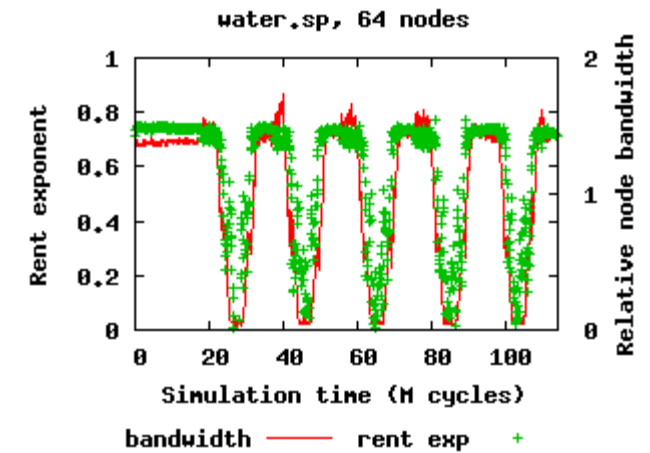
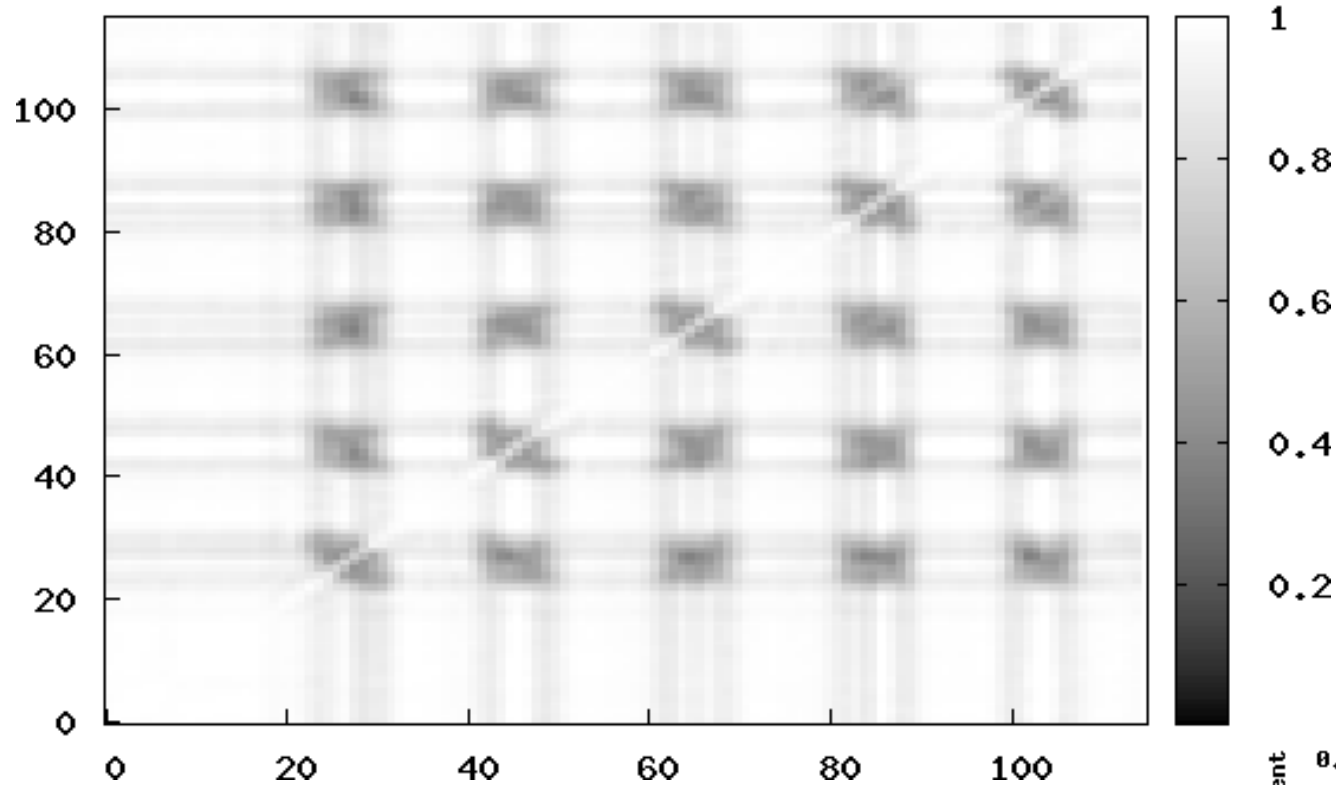
Similarity matrix: FFT

fft4M, 16 nodes



Similarity matrix: Water

water.sp, 64 nodes

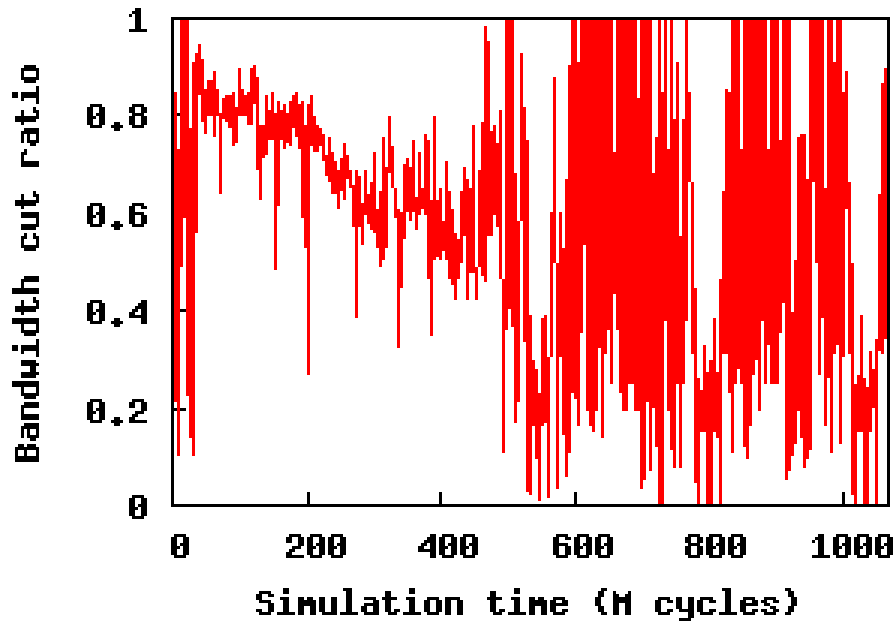


Suitability of a single placement

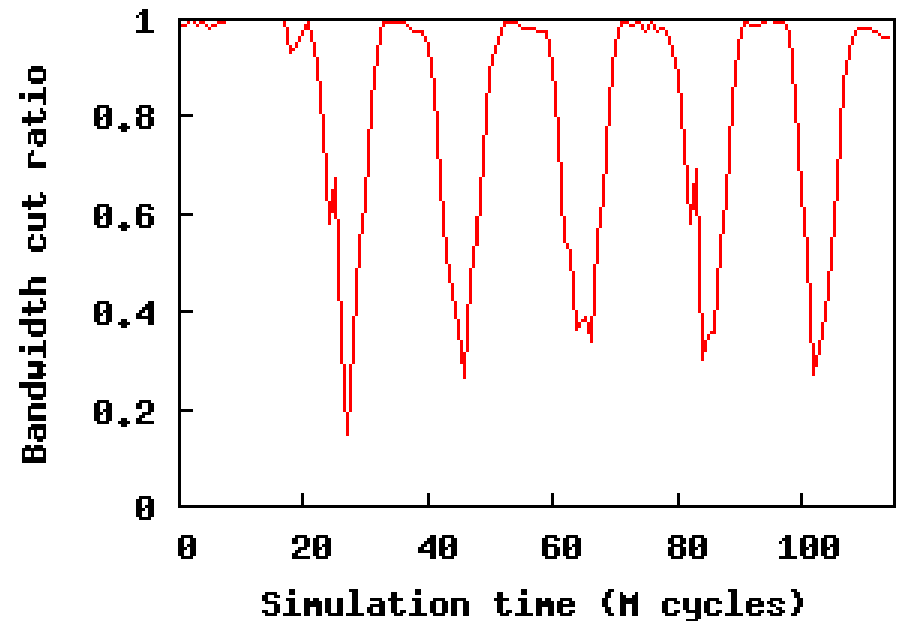
- Static network \rightarrow one single placement
- How suitable is this placement through time?
- Suitability measure: based on partitionings (as are placements)
- Optimal partitioning for traffic[X]:
part[X], cutting a bandwidth cut[X,X].
- Suitability of partitioning P:
 $\text{cut}[X,X] / \text{cut}[P,X]$

Suitability of a single placement

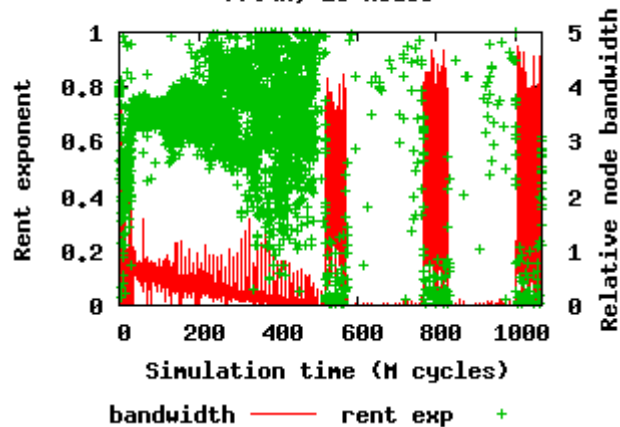
fft4M, 16 nodes



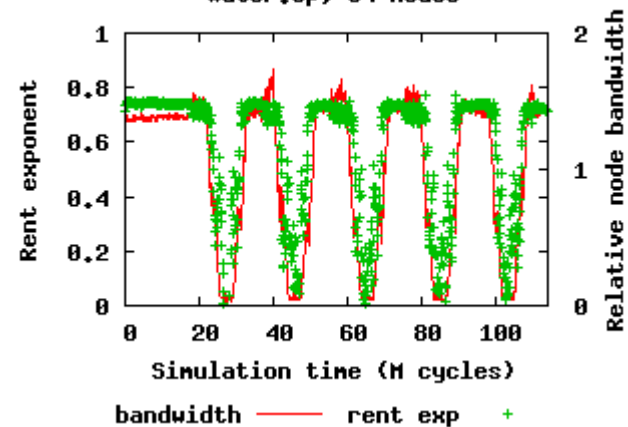
water.sp, 64 nodes



fft4M, 16 nodes



water.sp, 64 nodes



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Conclusions

- Measuring Rent exponents:
 - small number of nodes: difficult to measure, lots of noise
 - shared-memory: implicit communication, lots of non-essential communication → better/other results with message-passing?
- Still, difference in locality is visible, can be traced back to the benchmark's algorithm
- Time-variant communication!
- Rent's Rule (partitioning) is helpful to study communication behavior



Thank you!

