

# Is Probabilistic Congestion Estimation Worthwhile?

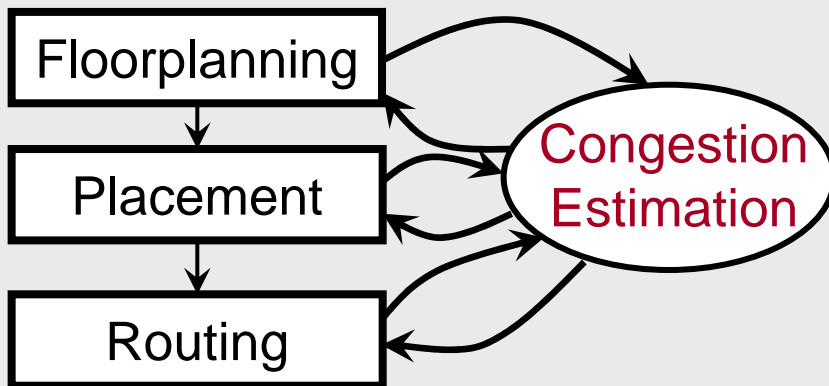
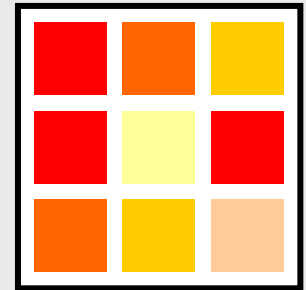
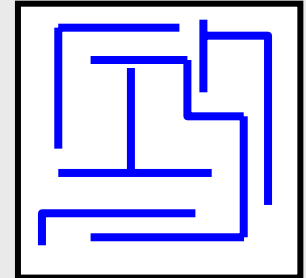
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# ❖ Introduction

- ◆ What is (Global) Routing?
  - Matching routing supply and demand.
- ◆ What is Congestion Estimation?
  - *Guessing* where matching is difficult!
- ◆ Why Congestion Estimation?
  - Prevent routability problems later on!



- ◆ Traditionally: Global Routing
  - Accurate / slow
- ◆ 2001: Probabilistic Methods 
  - Fast / less accurate

# ❖ Outline

- ◆ Motivation
- ◆ Implementation of probabilistic method
- ◆ Implementation of new GR-based method 
- ◆ The benchmarks
- ◆ Estimation quality
- ◆ Results
- ◆ Discussion
- ◆ Conclusion

# ❖ Motivation

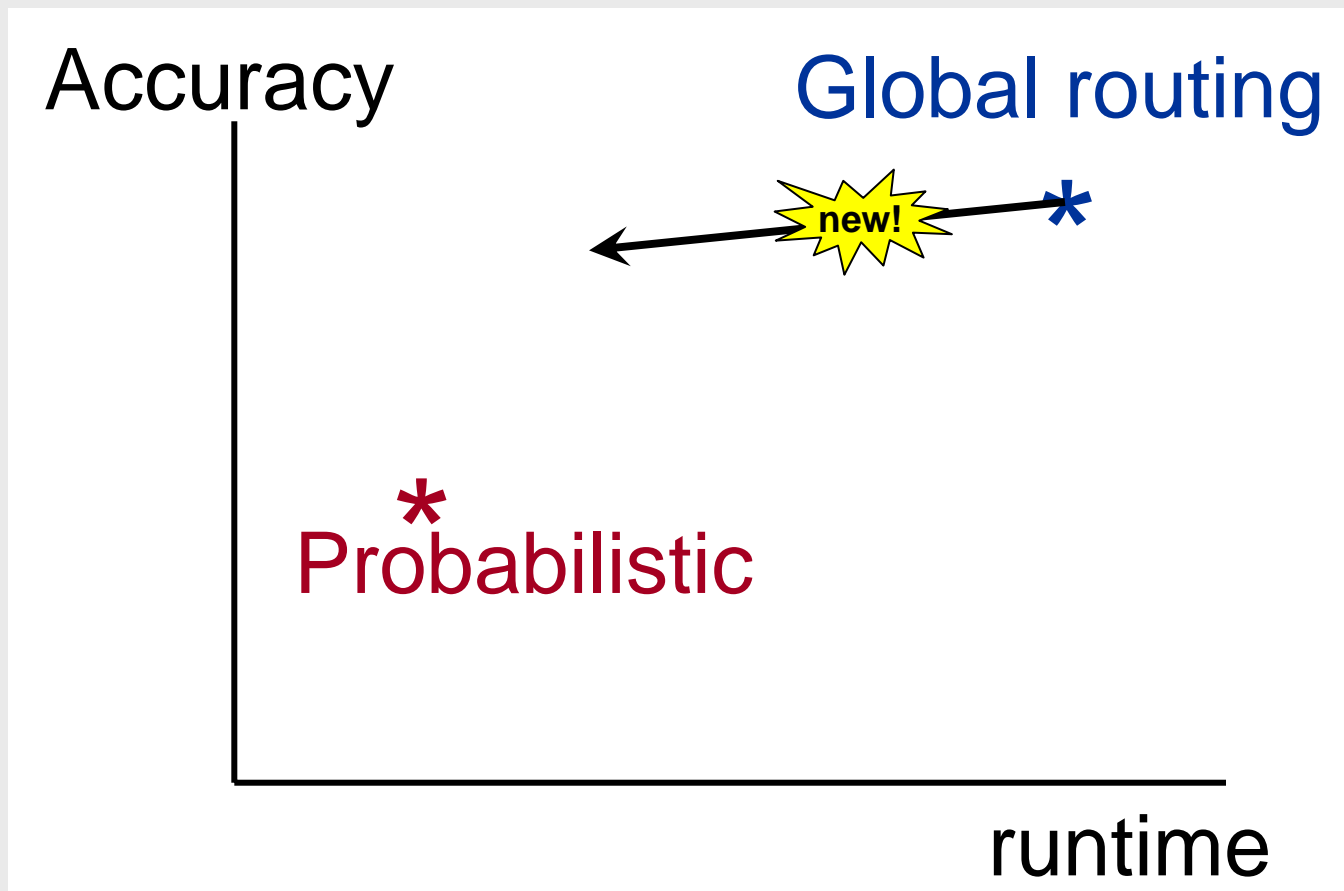
- ◆ Probabilistic methods **differ** in some of their observations.
  - Biased towards **tool** or **benchmarks!**
- ◆ Fundamentally, there is **little awareness of congestion** during probabilistic analysis.
  - Re-spreading afterwards.

Global routing is tuned towards  
minimizing congestion.

What if we tune it towards  
congestion estimation?



# ❖ Motivation



# ❖ Implementation of pce (Westra\*)

```
create maps h,v  
create hashmap stamplib  
break up nets with RMST
```

```
foreach wire w do
```

```
  if stamplib contains w do
```

```
    s ← get_stamp(stamplib,w)
```

```
    stamp_maps(h,v, w,s)
```

```
  else do
```

```
    s ← new_stamp(w)
```

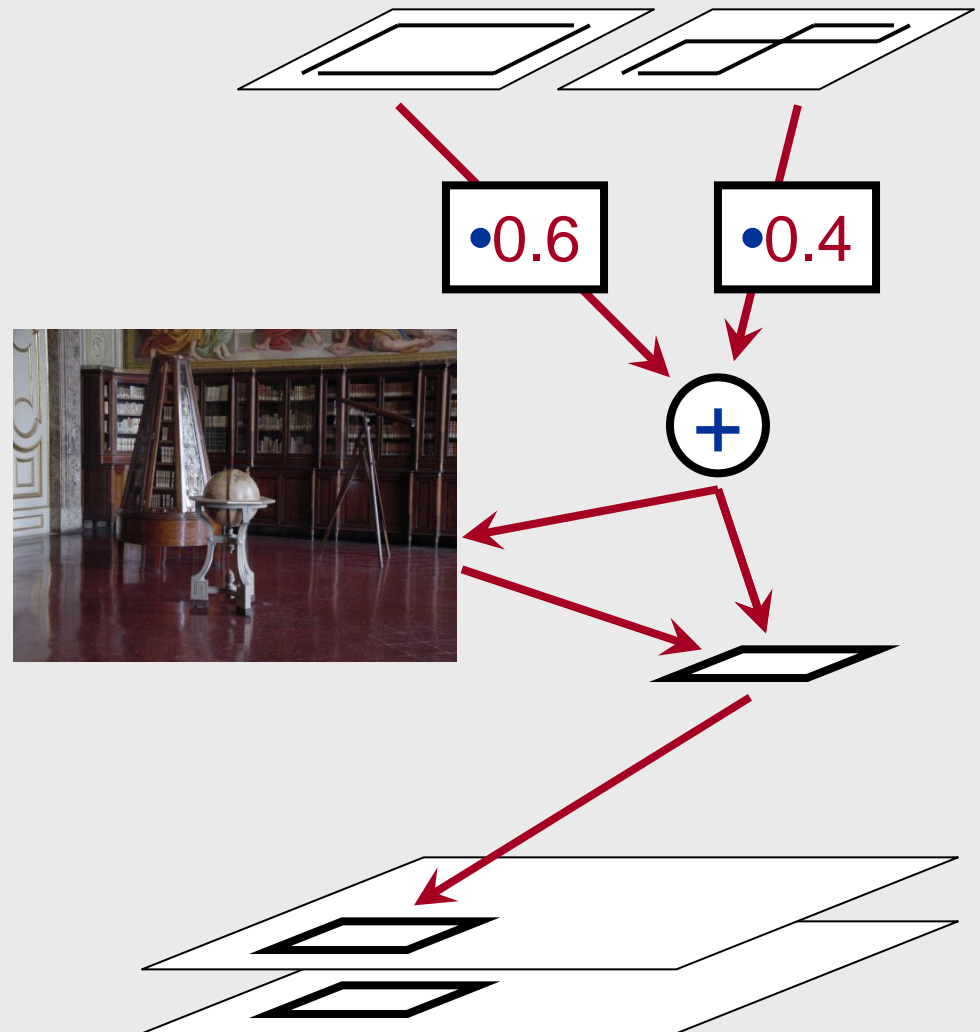
```
    put_stamp(stamplib,w,s)
```

```
    stamp_maps(h,v, w,s)
```

```
  end
```

```
end
```

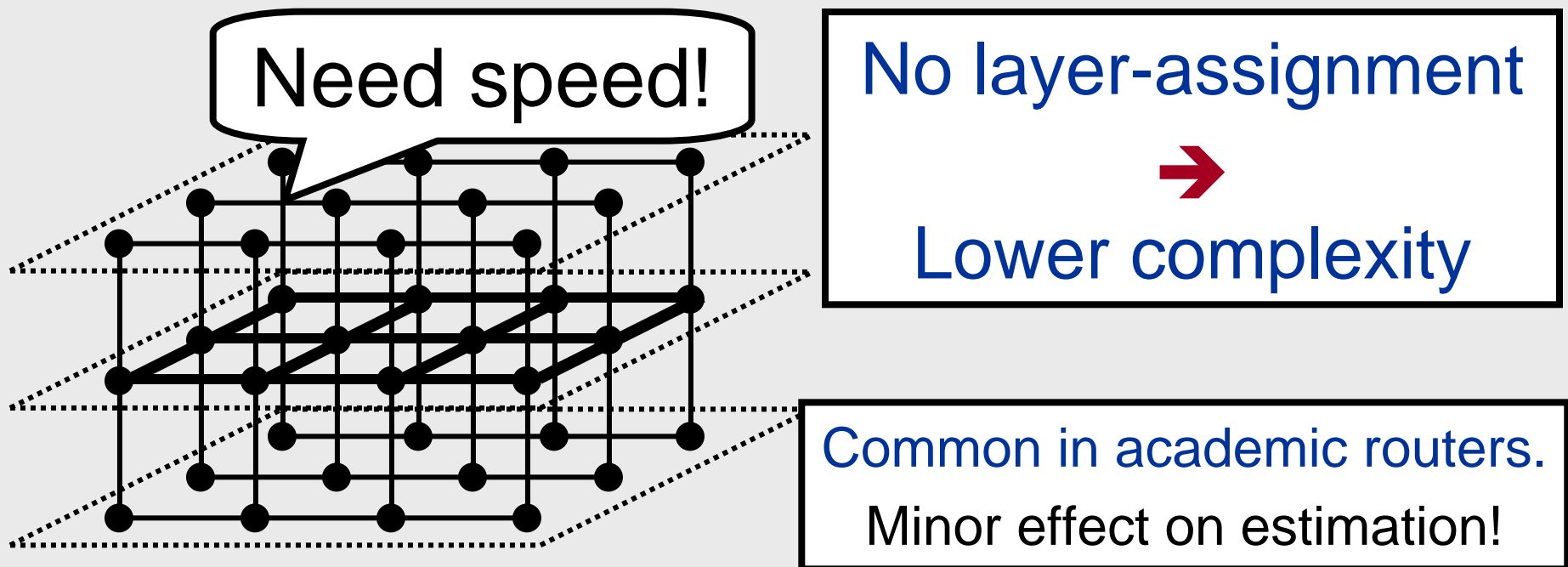
```
divide h and v by capacities
```



\*J. Westra *et al.*, "Probabilistic Congestion Prediction", ispd04  
SLIP05 San Francisco

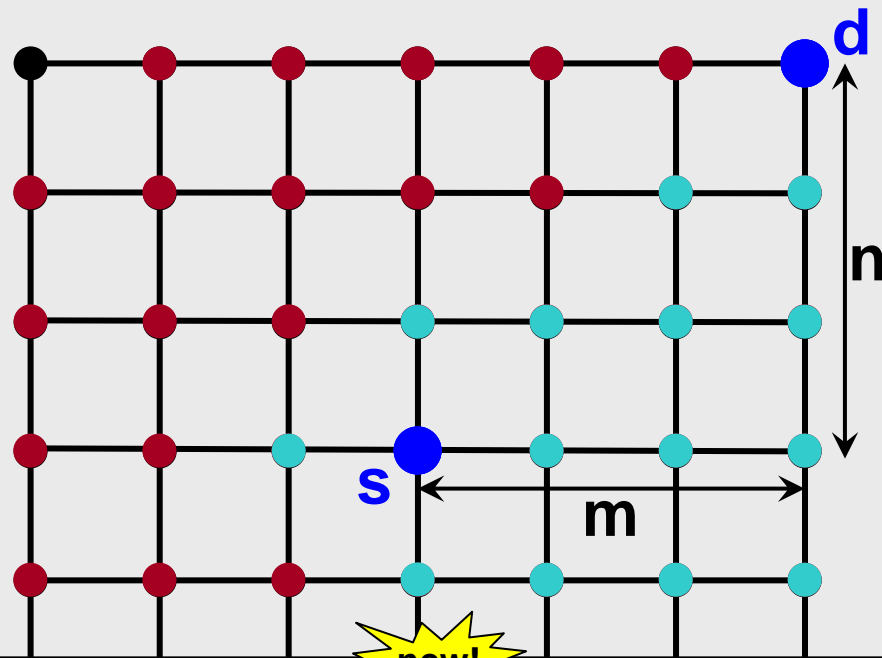
# ❖ Fast degenerate global routing

- ◆ FaDGloR<sup>new!</sup> is a maze router tuned towards speed and congestion estimation.
- ◆ Degenerate global routing model



# ❖ Fast degenerate global routing

- ◆ Common in GR: Dijkstra's algorithm.
- ◆ Faster:  $A^*$  algorithm.



## ◆ Visited nodes:

- Dijkstra:  $O((m+n)^2)$
- $A^*$ :  $O(m+n)$

If uncongested!

◆ FaDGloR: Smart  $A^*$  implementation!



# ❖ Fast degenerate global routing

## ◆ GR: ripup-and-reroute (R&R).

- First, allow **overflow**, **then spread** congestion
- Time consuming

Wire is routed **X** times!

## ◆ ~~FaDGloR: R&R.~~

- Make the **tradeoff** between wire length and congestion in one shot!

## ◆ FaDGloR: 2 step strategy:

- First route with **no overflow** → unrouted wires.
- Route remaining wires for **minimum overflow**.

# ❖ Fast degenerate global routing

## ◆ Wire ordering: common in GR: shortest first

- Literature: little effect

R&R!

- Experiments:

- Longest first → Less unrouted distance after 1<sup>st</sup> phase!
- Shortest first → Less overflow after 2<sup>nd</sup> phase!

contradiction?

Detouring!

new!

- ◆ **FaDGloR**: Shortest wires first.

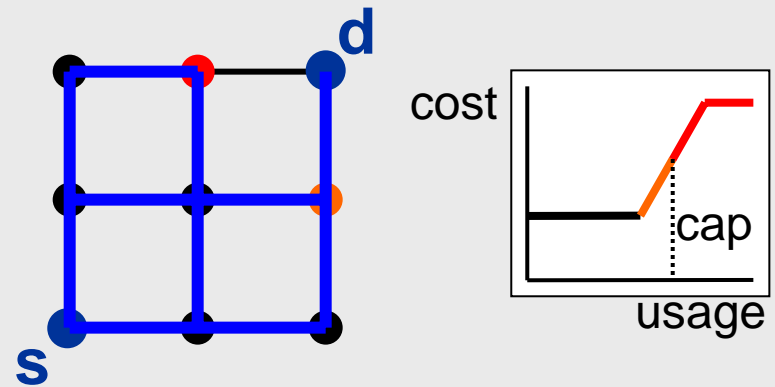
# ❖ Fast degenerate global routing

- ◆ GR: Cost-functions: path segments have cost
  - Penalize distance and local congestion
  - Increase cost of congestion during R&R

◆ CFs effective during R&R.

- FaDGloR: no congestion information available!

◆ CFs → more nodes visited!



Longest wires last  
→  
problem worse!

- new!**
- ◆ FaDGloR:
    - ~~Cost-function~~
    - ~~Detour-bounding~~

# ❖ The benchmarks

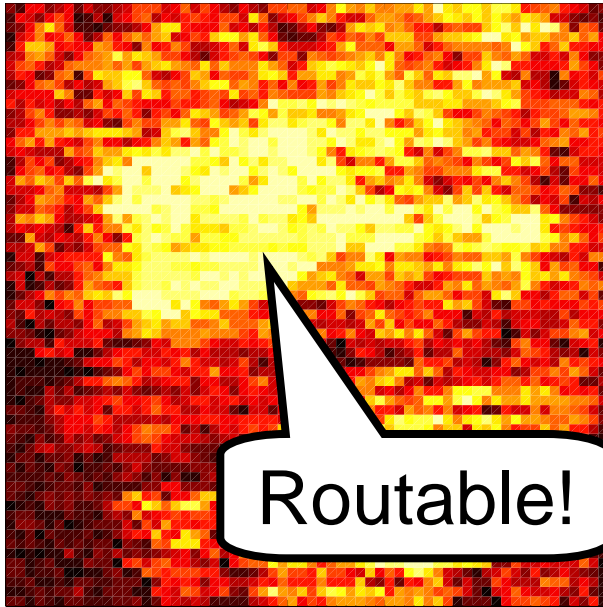
- ◆ The **Labyrinth** benchmarks
  - Only commonly used GR benchmarks (?)
  - From **real** designs.
- ◆ But they are all **difficult?!?!?!?**
  - In reality very **easy** to **impossible** designs!

Chip	Grid	Nets	wires	Chip	Grid	Nets	wires
ibm01	64x64	12k	27k	ibm06	128x64	33k	79k
ibm02	80x64	18k	53k	ibm07	192x64	44k	105k
ibm03	80x64	22k	44k	ibm08	192x64	48k	128k
ibm04	96x64	26k	52k	ibm09	256x64	50k	124k
<b>ibm05</b>	<b>128x64</b>	<b>28k</b>	<b>90k</b>	ibm10	256x64	64k	175k

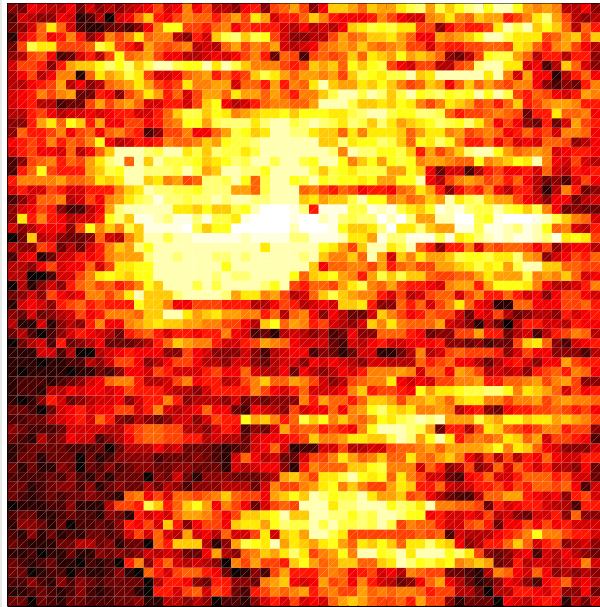


◆ **FaDGloR**: Added capacity -5...+10.

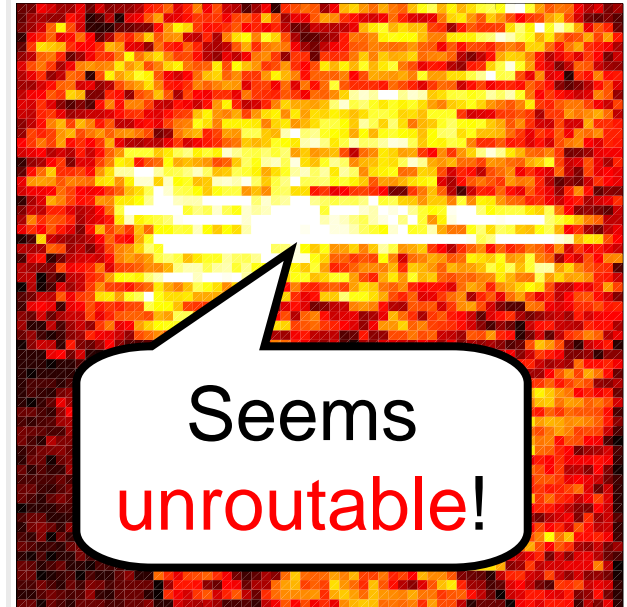
# ❖ Estimation quality



Labyrinth



FaDGloR

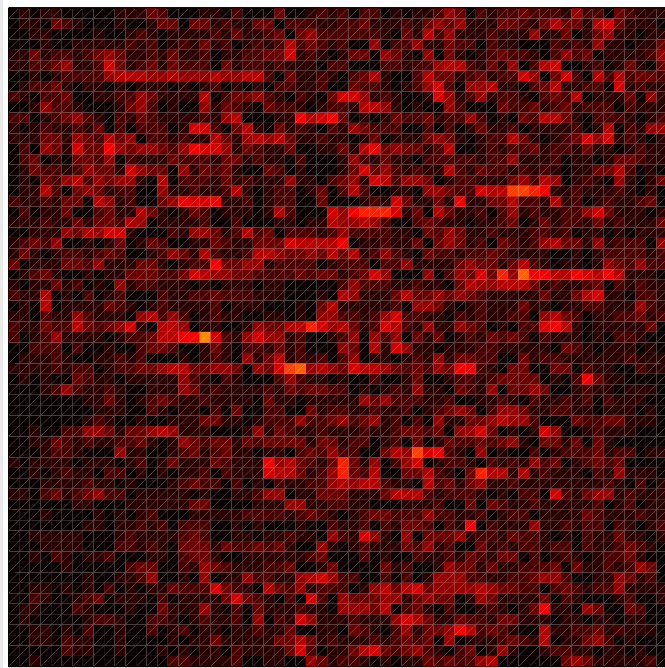


pce

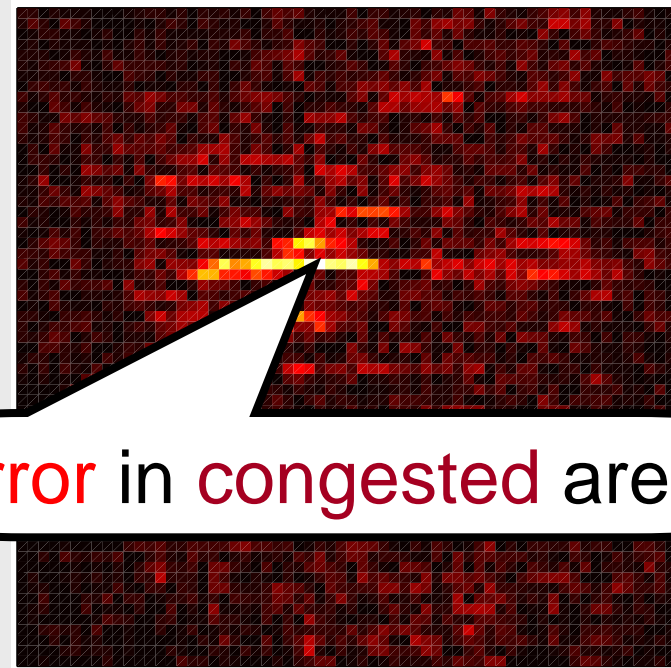
- ◆ Interested in 'wrongly congested'.
  - $c(i,j) > 1.1 \wedge C(i,j) \leq 1.1$
- ◆ Less interested in 'wrongly uncongested'.
  - $c(i,j) < 0.9 \wedge C(i,j) \geq 0.9$

# ❖ Estimation quality

## Estimation errors



FaDGloR



Error in congested area!

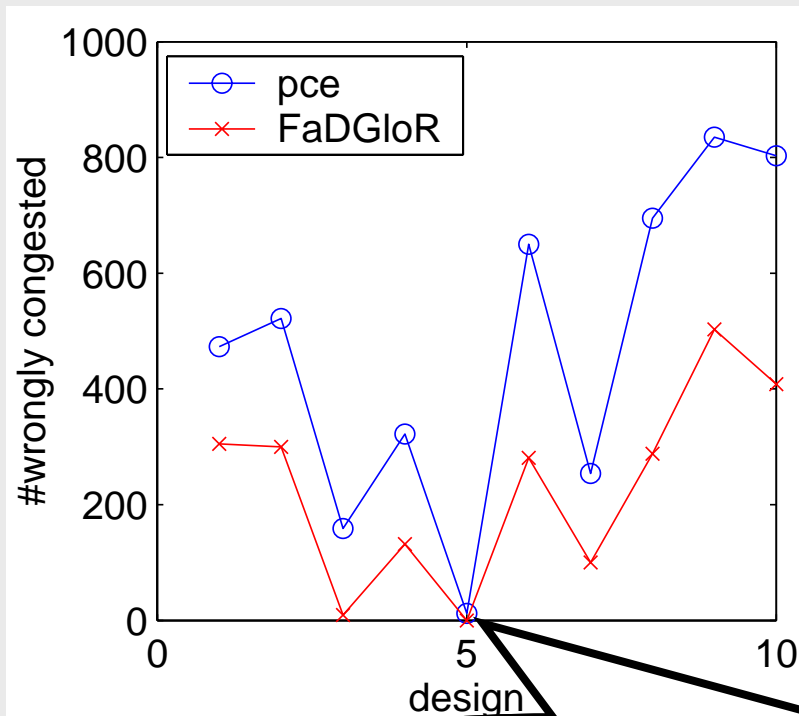
pce

new!

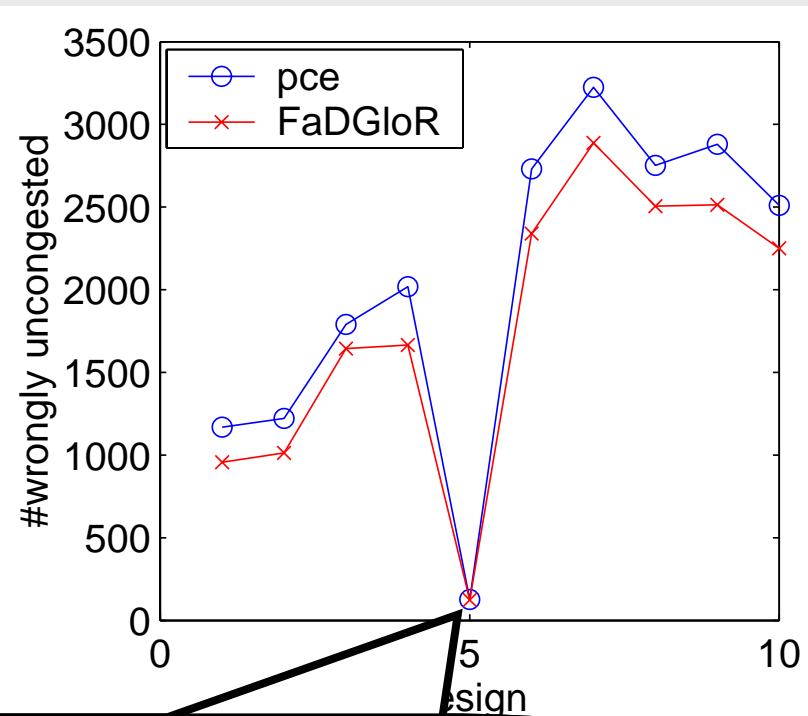
- ◆ **FaDGloR:** Errors more noise-like!
  - Error independent of congestion level.

# ❖ Results

## Wrongly congested



## Wrongly uncongested

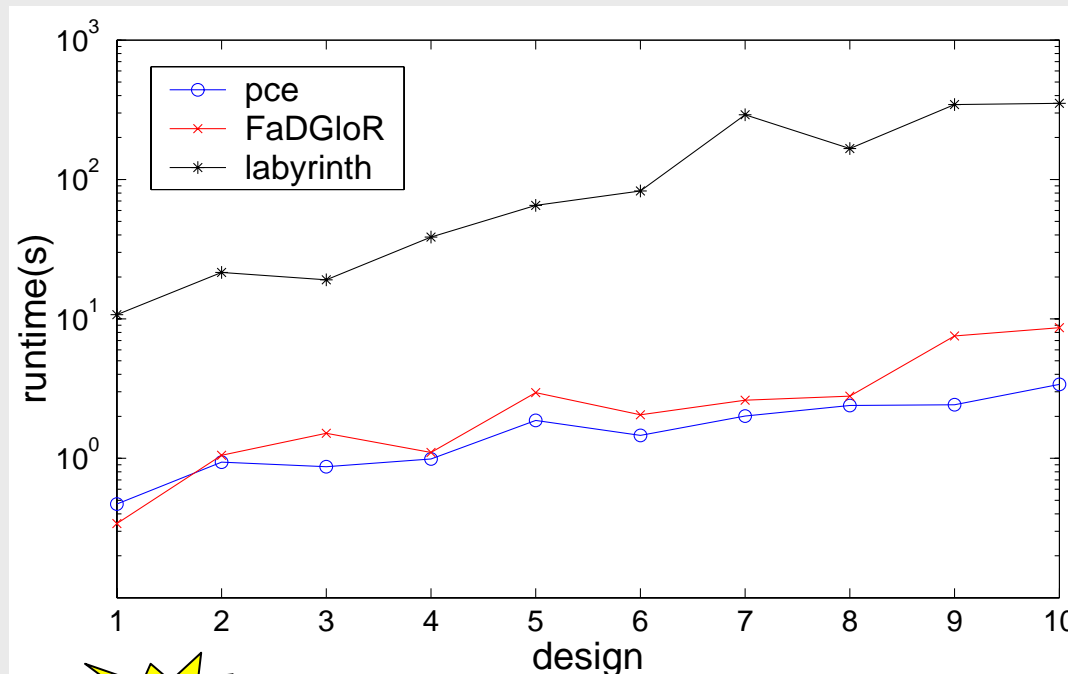


**new!**

- ◆ **FaDGloR: Half the wrongly congested!**
- ◆ **Slightly less wrongly uncongested!**

# ❖ Runtimes

- ◆ Capacity dependence → First routable
  - Still difficult designs!

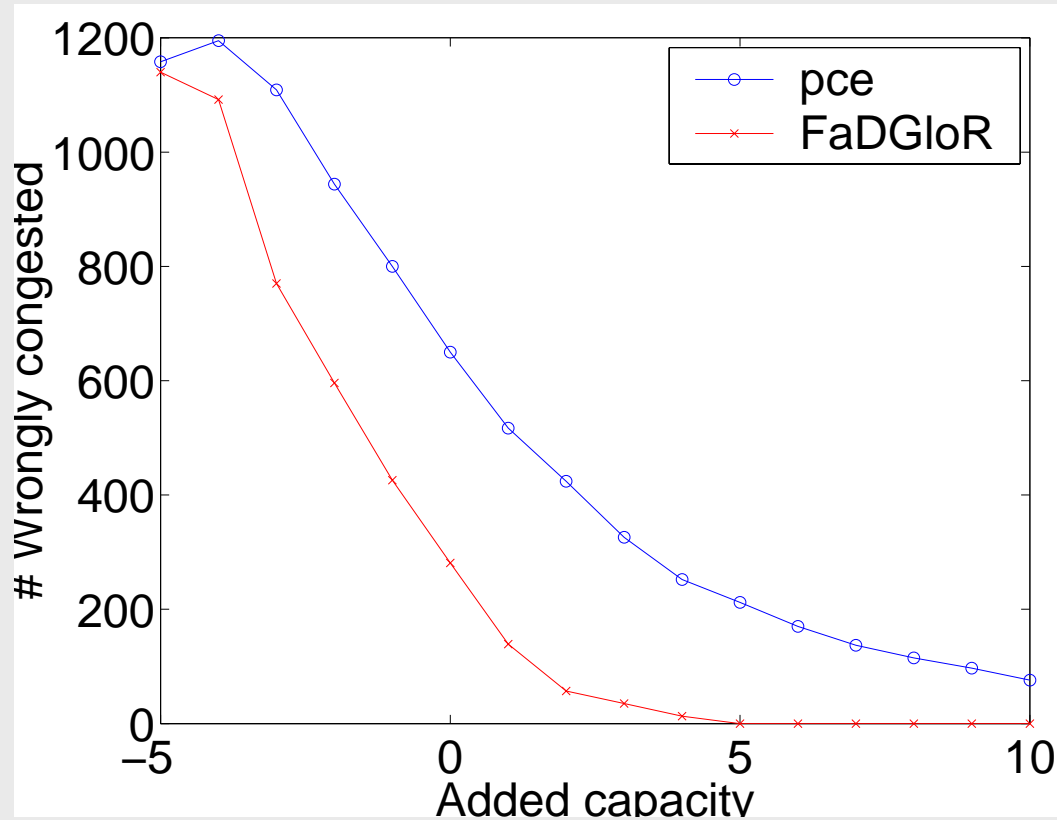


**new!**

- ◆ **FaDGloR**: Only slightly slower than pce, but order of magnitude faster than Labyrinth!

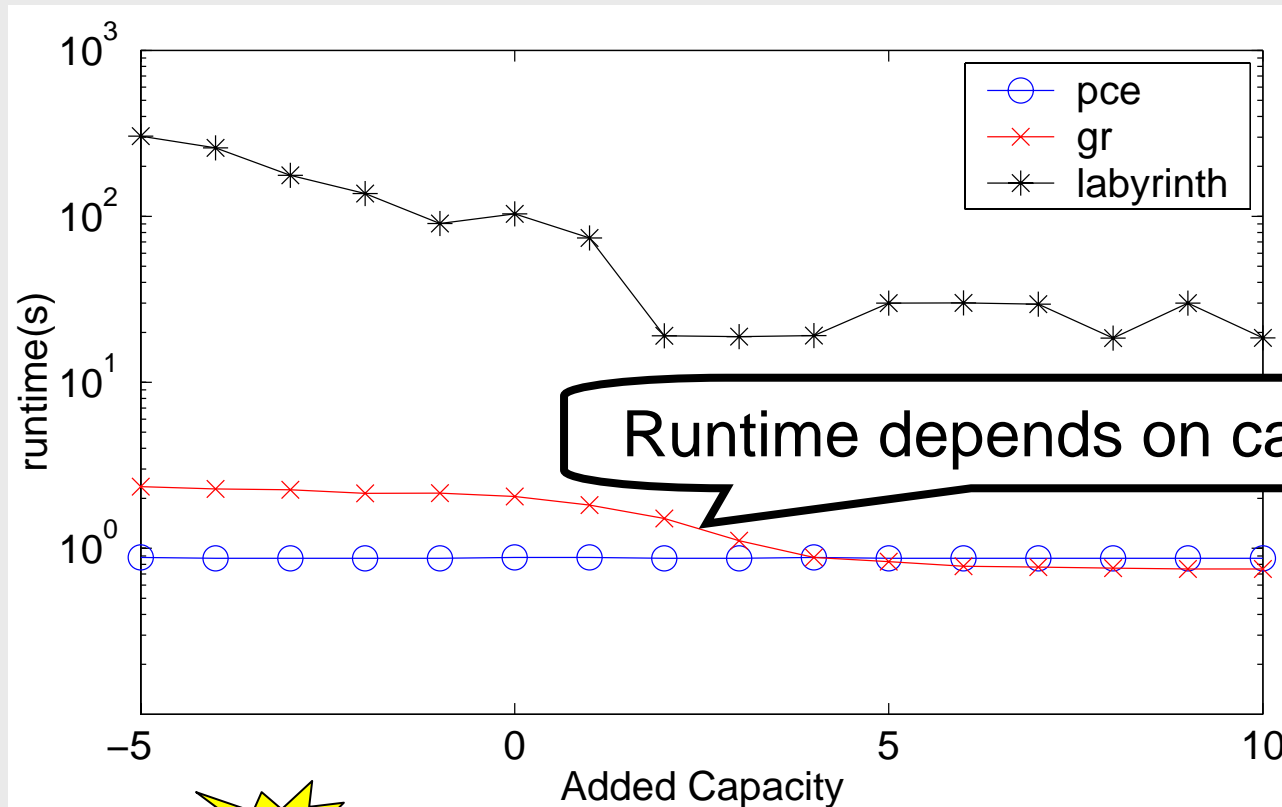


# ❖ Adding capacity - Quality



- ◆ **FaDGloR**: Much less wrongly congested at all capacities!

# ❖ Adding capacity - Runtimes



- ◆ **FaDGloR**: **Faster with more capacity!**
  - **High cap**: sometimes faster, sometimes slower than **pce**!

# ❖ Discussion

- ◆ **FaDGloR** <sup>new!</sup> much better on 'wrongly congested'
  - Slightly congestion-driven
  - Detouring allowed
- ◆ **FaDGloR** <sup>new!</sup> only slightly better on 'wrongly uncongested'
  - Labyrinth detours much more → more congestion
  - uncongested → many legal realizations




# ❖ Discussion

- ◆ **FaDGloR** <sup>new!</sup> only slightly slower than pce
  - pce: stamps → m·n entries copied
  - FaDGloR: uncongested → O(m+n) nodes visited
- ◆ Is pce a fast implementation?

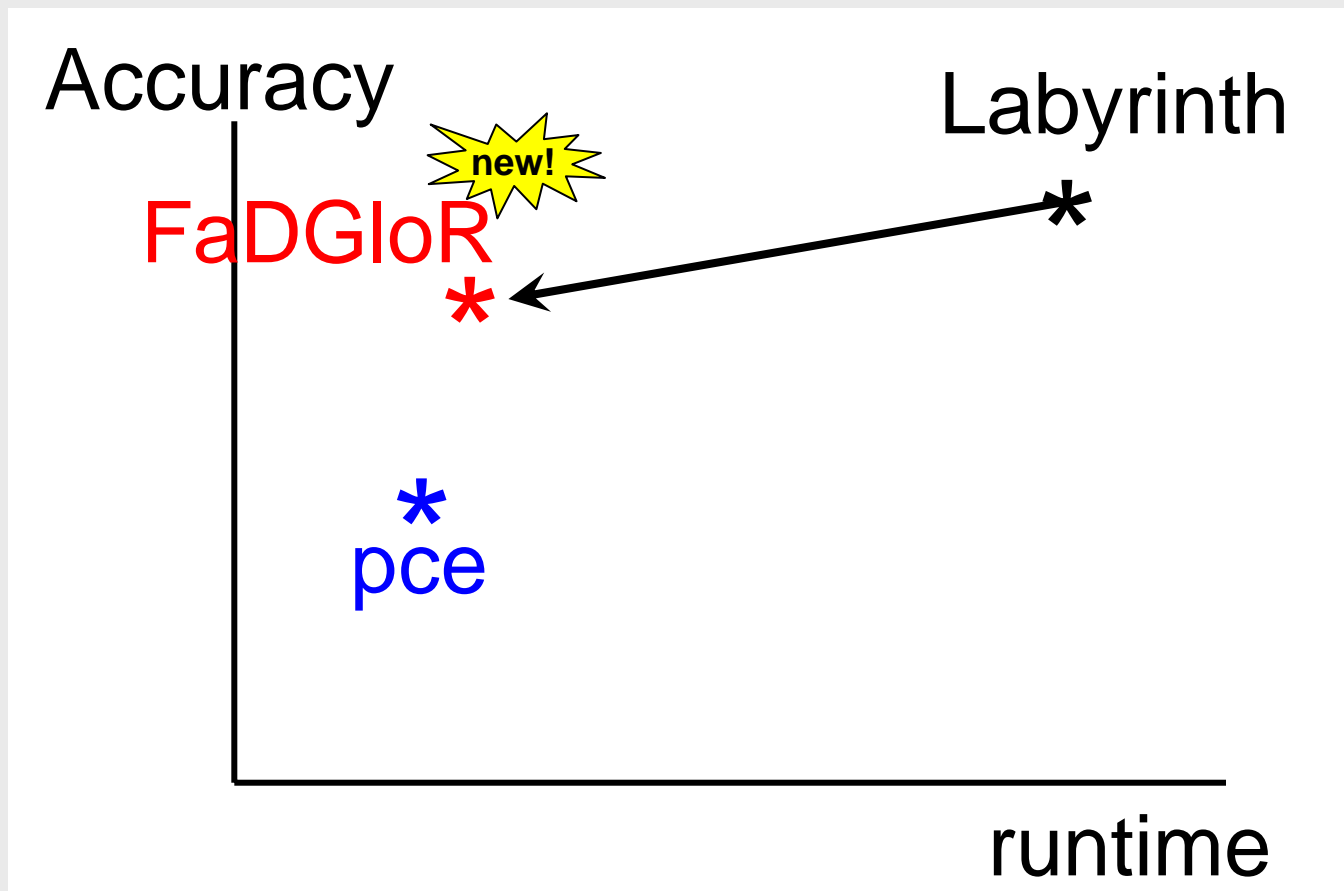
Method	Runtime per net		cpu
		1 <sup>st</sup> routable	
Lou	250us		???
Kahng	330us		2.4 GHz
Pce I	50us		1.0 GHz
Pce II	14us		1.4 GHz
FaDGloR I	88us	77us	1.0 GHz
FaDGloR II	40us	35us	1.4 GHz

- Ongoing improvements... **at the price of runtime!**

# ❖ Discussion

- ◆ **FaDGloR**:  one shot solution
  - Highly tuned: interplay between 2-step strategy, detour-bounding and wire ordering
- ◆ **FaDGloR**  probably better with blockages than pce
- ◆ **FaDGloR**  allows refinement
  - R&R certain windows
  - **FaDGloR** results are **seeds** for other applications

# ❖ Discussion



## ❖ Conclusion

Is Probabilistic Congestion Estimation Worthwhile?

NO!

Global routing-based methods  
more flexible and promising!

# Thank you!

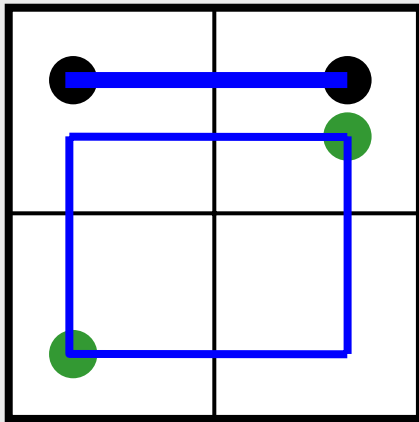
## Questions?



# ❖ Problem with stamping

- ◆ Stamping: no awareness of congestion!

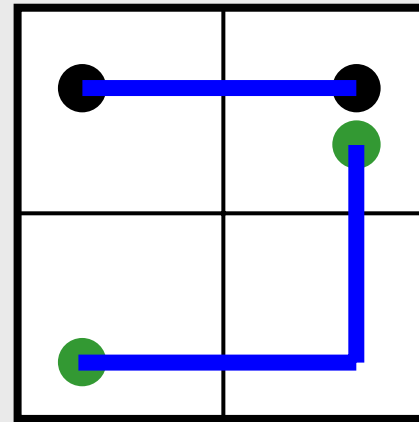
Capacity = 1



$1\frac{1}{2}$

$\frac{1}{2}$

pce



1

1

FaDGloR

- ◆ Solution: re-spreading afterwards → slow!

# ❖ Motivation – previous work

- ◆ Lou *et al.*: “Estimating Routing Congestion using Probabilistic Analysis” ispd 2001
  - Tiles have fixed capacity
  - Spread wire over detour-free paths with equal probability
- ◆ Kahng *et al.*: “Accurate Pseudo-Constructive Wirelength Estimation” ispd 2003
  - Us
  - Co
- ◆ West *et al.*: “Accurate Pseudo-Constructive Wirelength Estimation” ispd 2003
  - L- and Z-shapes only
  - No detours

**Not consistent!**

**Biased towards tool  
and/or benchmarks!**